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# (12) United States Patent

NepomucenoLeung et al.

# (54) VIDEO REFRESH ADAPTATION ALGORITHMS RESPONSIVE TO ERROR FEEDBACK

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- (51) Int. Cl.

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  H04N 19/507 (2014.01)

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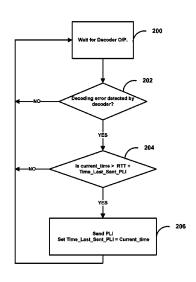
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# (57) ABSTRACT

In general, this disclosure relates to techniques for adaptively determining a rate for intra-coding segments (e.g., macro-blocks) of video data within a sequence of video units (e.g., frames) in response to the receipt of error feedback that is provided by a video decoder. One example method includes receiving feedback from a video decoder that indicates an error associated with coded video data, and, upon receipt of the feedback, determining a rate for intracoding a specified number of segments of video data for each video unit within a sequence of video units, wherein an intra-coding of segments of video data within one video unit is independent of a coding of segments of video data within other video units within the sequence.

# 34 Claims, 10 Drawing Sheets

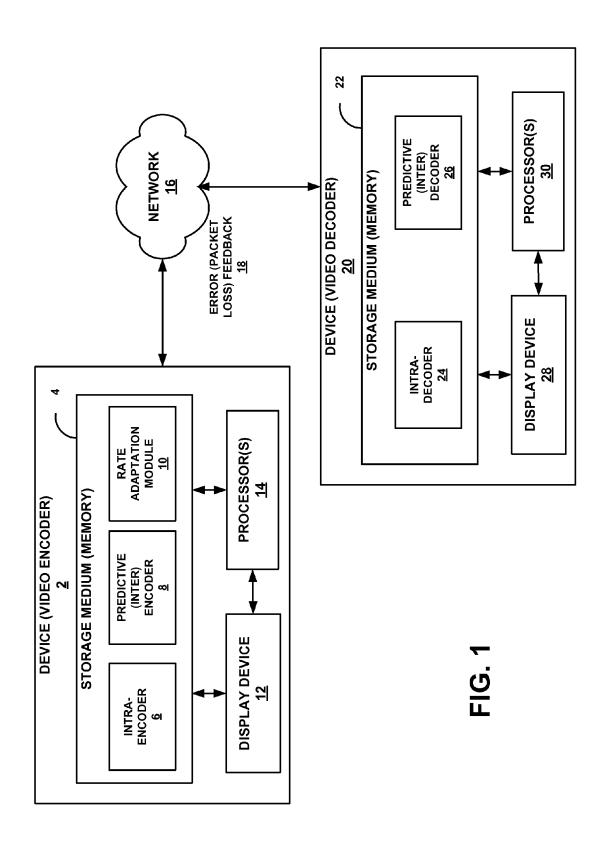


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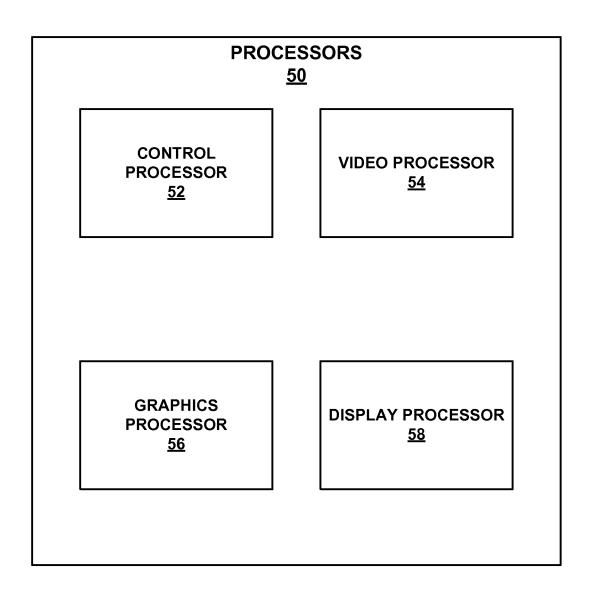


FIG. 2

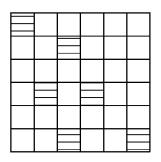


FIG. 3A

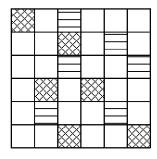


FIG. 3B

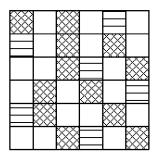


FIG. 3C

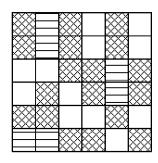


FIG. 3D

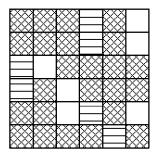


FIG. 3E

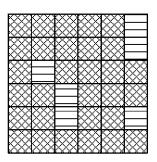
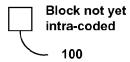
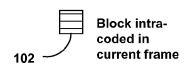
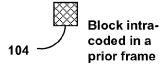


FIG. 3F







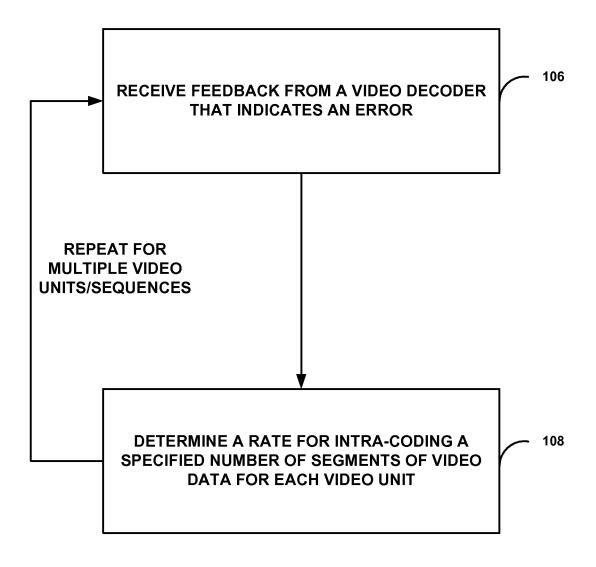
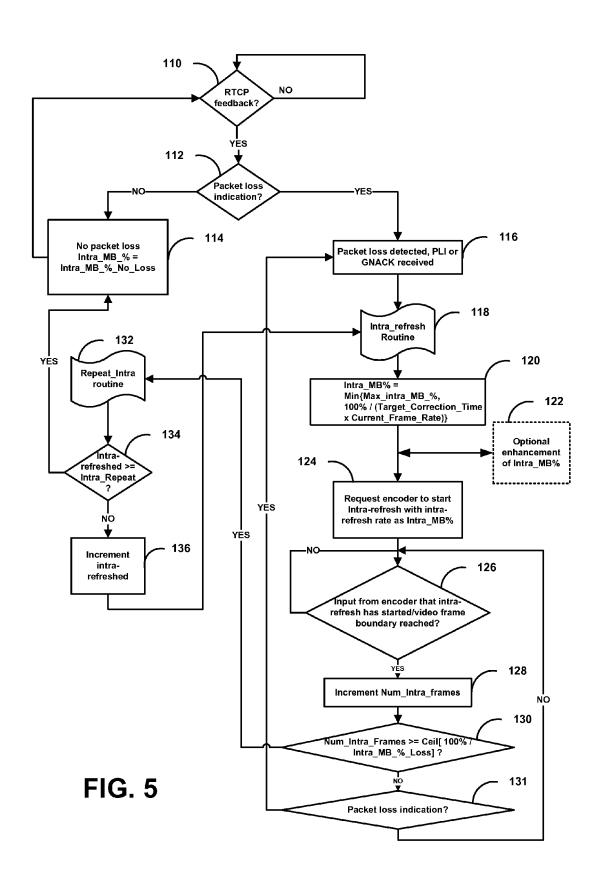
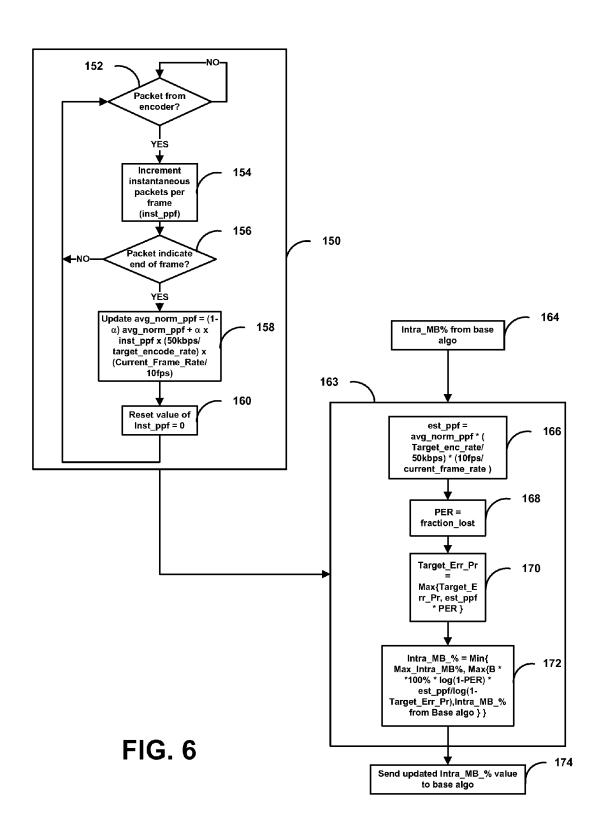


FIG. 4





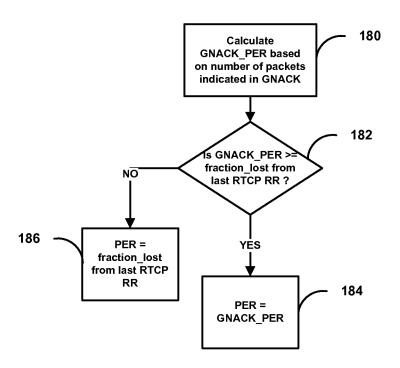


FIG. 7

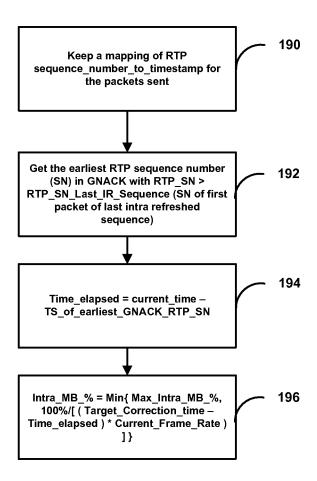


FIG. 8

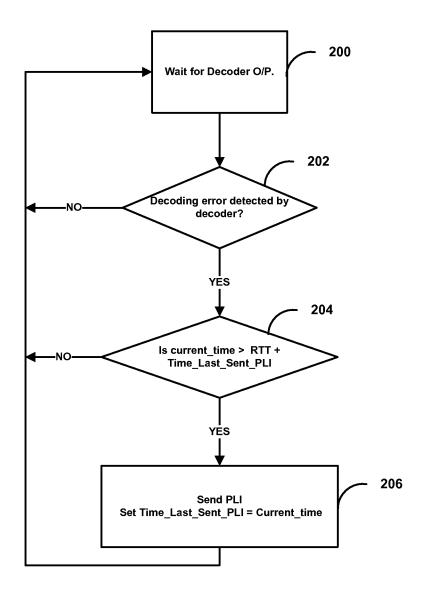
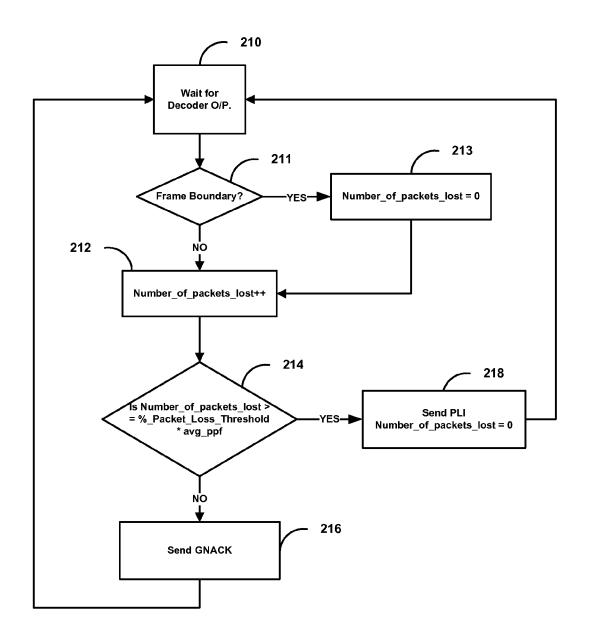


FIG. 9



**FIG. 10** 

# VIDEO REFRESH ADAPTATION ALGORITHMS RESPONSIVE TO ERROR FEEDBACK

### RELATED APPLICATION

This application is a divisional application of U.S. patent application Ser. No. 12/349,386, filed Jan. 6, 2009, which claims the benefit of U.S. Provisional Application No. 61/043,005, filed on Apr. 7, 2008, the entire contents of both of which are incorporated herein by reference.

### TECHNICAL FIELD

This disclosure relates to encoding and/or decoding of <sup>15</sup> video data.

## BACKGROUND

Video coding and compression techniques may achieve <sup>20</sup> efficient compression by reducing both temporal redundancies between video frames in a frame sequence and also spatial redundancies within a video frame. Examples of video coding and compression techniques are described by the ITU-T H.261, H.263, Motion Picture Experts Group <sup>25</sup> (MPEG) 1, MPEG2, and MPEG4 standards, as well as the ITU-T H.264 standard and its counterpart, ISO/IEC MPEG-4, Part 10, i.e., Advanced Video Coding (AVC).

An intra-coded frame is self-contained and reduces spatial redundancies within a video frame. An inter-coded frame, 30 however, is predicted via motion compensation from other coded frames (e.g., previously coded frames) to reduce temporal redundancies. Inter-coding may achieve high compression efficiency. However, the inter-dependency between frames makes the coded bit-stream for inter-coded frames more susceptible to propagating channel errors. Errors introduced in the compressed bit-stream may result in errors in the video frames reconstructed by a video decoder.

Intra-coding may help minimize the possibility of propagating channel errors, because the coding of video data in these frames is independent of the coding of video data in previous or subsequent frames. However, intra-coded frames may comprise many times more coded data than inter-coded frames, and may therefore cause transmission bandwidth issues. As a result, there may be a degradation of quality on the decoded video output side, particularly when decoding is performed by real-time, or streaming, video applications. Thus, in some cases, video coding techniques may only intra-code individual segments (e.g., macroblocks) of data within individual frames, rather than intra-coding all of the frame data.

Any given macro-block within an inter-coded frame (i.e., Predicted frame (P-frame) or Bi-directionally predicted frame (B-frame)) may be intra-coded as an intra macro-block. Similar to intra-coded frames, an intra-coded macro-55 block is coded independently of data from another (e.g., previously) coded frame. The method of forcing macro-blocks to be intra-coded may be referred to as encoder intra-refresh.

# SUMMARY

In general, this disclosure relates to techniques for adaptively determining a rate for refreshing segments (e.g., macro-blocks) of video data within a sequence of video units 65 (e.g., frames) in response to the receipt of error feedback provided by a video decoder. This error feedback may

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comprise packet loss feedback, indicating that the video decoder failed to receive packets within one or more frames of a frame sequence.

In some cases, the error feedback comprises a Picture Loss Indication (PLI) message and/or a General Negative Acknowledgment (GNACK) message that is sent by the video decoder. A PLI message indicates a loss of an undefined amount of coded video data belonging to one or more pictures, while a GNACK message may identify specific packets of coded video data that have been lost.

An error feedback algorithm may be implemented to set a macro-block intra-refresh rate by the video encoder based upon the error feedback that is received from the video decoder. The video encoder may then use this intra-refresh rate to intra-code one or more macro-blocks in each frame of a frame sequence.

In one aspect, a method comprises receiving feedback from a video decoder that indicates an error associated with coded video data, and, upon receipt of the feedback, determining a rate for intra-coding a specified number of segments of video data for each video unit within a sequence of video units, wherein an intra-coding of segments of video data within one video unit is independent of a coding of segments of video data within other video units within the sequence.

In one aspect, a computer-readable medium contains instructions for causing one or more processors to receive feedback from a video decoder that indicates an error associated with coded video data, and, upon receipt of the feedback, to determine a rate for intra-coding a specified number of segments of video data for each video unit within a sequence of video units, wherein an intra-coding of segments of video data within one video unit is independent of a coding of segments of video data within other video units within the sequence.

In one aspect, a device comprises a storage device and one or more processors coupled to the storage device and configured to receive feedback from a video decoder that indicates an error associated with coded video data, and, upon receipt of the feedback, to determine a rate for intracoding a specified number of segments of video data for each video unit within a sequence of video units, wherein an intra-coding of segments of video data within one video unit is independent of a coding of segments of video data within other video units within the sequence.

In one aspect, a method comprises detecting, by a video decoder, at least one error associated with coded video data sent from a video encoder, determining a round-trip travel time for data communication between the video decoder and the video encoder, determining a last-sent time at which the video decoder last sent the video encoder a feedback message indicating a previous error that was detected by the video decoder, comparing a current time to a sum of the last-sent time and the round-trip travel time, and, based upon the comparing, determining whether to send the video encoder a second feedback message indicating the at least one error detected by the video decoder.

In one aspect, a device comprises a storage device and one or more processors coupled to the storage device. The one or more processors are configured to detect at least one error associated with coded video data sent from a video encoder, determine a round-trip travel time for data communication between the device and the video encoder, determine a last-sent time at which the device last sent the video encoder a feedback message indicating a previous error that was detected by the device, and compare a current time to a sum of the last-sent time and the round-trip travel time. Based

upon the comparison, the one or more processors are further configured to determine whether to send the video encoder a second feedback message indicating the at least one error detected by the device.

The techniques described in this disclosure may be implemented in hardware, software, firmware, or any combination thereof. If implemented in software, the software may be executed in a processor, which may refer to one or more processors, such as a microprocessor, application specific integrated circuit (ASIC), field programmable gate array (FPGA), or digital signal processor (DSP), or other equivalent integrated or discrete logic circuitry. Software comprising instructions to execute the techniques may be initially stored in a computer-readable medium and loaded and executed by a processor.

Accordingly, this disclosure also contemplates computerreadable media comprising instructions to cause a processor to perform any of a variety of techniques as described in this disclosure. In some cases, the computer-readable medium may form part of a computer program product, which may be sold to manufacturers and/or used in a device. The 20 computer program product may include the computer-readable medium, and in some cases, may also include packaging materials.

The details of one or more aspects are set forth in the accompanying drawings and the description below. Other features, objects, and advantages will be apparent from the description and drawings, and from the claims.

### BRIEF DESCRIPTION OF DRAWINGS

FIG. 1 is a block diagram illustrating a video encoder that is communicatively coupled to a video decoder via a net-

FIG. 2 is a block diagram illustrating example processors that may be included within the video encoder and/or video decoder shown in FIG. 1.

FIGS. 3A-3F are diagrams illustrating a frame sequence, where a certain number of blocks are intra-coded within each frame of the sequence.

FIG. 4 is a flow diagram illustrating a method that may be performed by a video encoder to determine a rate for 40 intra-coding segments of video data based upon received error feedback.

FIG. 5 is a flow diagram illustrating a method that may be performed by a video encoder to determine an intra-refresh rate for intra-coding blocks of video data based upon 45 received packet loss information.

FIG. 6 is a flow diagram of a method that may be performed by a video encoder to obtain a modified value of the intra-refresh rate determined by the method of FIG. 5.

FIG. 7 is a flow diagram of a method that may be performed by a video encoder to calculate a packet error rate upon receipt of a General Negative Acknowledgment (GNACK) message.

FIG. 8 is a flow diagram illustrating a method that may be performed by a video encoder to determine a modified that includes packet sequence numbers.

FIG. 9 is a flow diagram illustrating a method that may be performed by a video decoder to determine whether to send a PLI (picture loss indication) message to a video encoder.

FIG. 10 is a flow diagram of a method that may be 60 performed by a video decoder to determine whether to send a GNACK or a PLI to a video encoder.

# DETAILED DESCRIPTION

FIG. 1 is a block diagram illustrating an example device (video encoder) 2 that is communicatively coupled to a

device (video decoder) 20 via network 16. Video encoder 2 and video decoder 20 may comprise stand-alone devices or may be part of larger systems. For example, each of video encoder 2 and video decoder 20 may comprise, or be part of, a wireless communication device (such as a wireless mobile handset), a digital camera, a video camera, a video telephone, a digital multimedia player, a personal digital assistant (PDA), a video game console, a personal computer or laptop device, or other video device. Video encoder 2 and video decoder 20 may also each be included within one or more integrated circuits, or chips, which may be used in some or all of the devices described above. Network 16 may, in some cases, comprise a wired or a wireless network. In one aspect, network 16 may be a private or a public (e.g., Internet) network.

For purpose of illustration and description only, device 2 is referred to as a video encoder, and device 20 is referred to as a video decoder. In the description below, various techniques provides details on exchanges of information between a device that performs encoding (e.g., compression) and a device that performs decoding (e.g., decompression), and thus device 2 (video encoder) and device 20 (video decoder) are shown as being communicatively coupled via network 16. It should be understood, however, that, in certain aspects, device 2 and/or device 20 may perform encoding and/or decoding operations. However, it is assumed in the subsequent description, for illustration purposes only, that device 2 performs at least video encoding operations, and is therefore referred to as a video encoder, while device 20 performs at least video decoding operations, and is therefore referred to as video decoder.

In some cases, video encoder 2 and video decoder 20 may each be capable of executing various applications, such as graphics applications, video applications, and/or other multi-media applications. For example, video encoder 2 and video decoder 20 may be used for graphics applications, video game applications, video playback applications, digital camera applications, instant messaging applications, video teleconferencing applications, mobile applications, or video streaming applications.

Video encoder 2 and video decoder 20 may each be capable of processing a variety of different data types and formats. For example, they may be able to process still image data, moving image (video) data, or other multimedia data. As will be described in more detail below, video encoder 2 is capable, during an intra-refresh procedure, of intra-coding one or more segments (e.g., blocks or macroblocks) or video data within one or more video units, such as video frames, within a sequence of video units to help reduce spatial redundancies within the sequence. Intracoded segments are segments that are coded independently of video data from a previously or subsequently coded video

Any given macro-block within an individual frame may intra-refresh rate based upon receipt of a GNACK message 55 be intra-coded as an intra macro-block. Similar to intracoded frames, an intra-coded macro-block is coded independently of data from another (e.g., previously) coded frame. The method of forcing macro-blocks to be intracoded may be referred to as encoder intra-refresh. The rate at which macro-blocks are intra-coded may be referred to as an intra-refresh rate. There may be various reasons for intra-refreshing macro-blocks, such as, for example, inverse discrete cosine transform (IDCT) mismatch control by video decoder 20, and error resilience.

> To control IDCT mismatch, such as within the context of the H.261 and H.263 block-based video coding standards, each macro-block location in an image may be intra-coded

at least once every pre-determined number of times that coefficients are transmitted for that macro-block. The intrarefresh of macro-blocks may limit the extent of error propagation due to DCT/IDCT mismatch.

To improve error resilience, selected macro-blocks may 5 be intra-coded to limit error propagation resulting from using corrupt macro-blocks that have been incorrectly reconstructed or concealed due to channel errors. These corrupt macro-blocks may, upon decoding, be visually objectionable. Furthermore, correctly decoded macro-blocks from subsequent frames that refer back to a corrupt macro-block for temporal prediction may also be visually objectionable. These types of artifacts may be more objectionable than the DCT/IDCT mismatch errors. As such, they may drive an intra-refresh strategy when communicating 15 data over error-prone channels.

Based upon error (packet loss) feedback 18 that is received from video decoder 20, video encoder 2 is capable of dynamically adjusting the rate of intra-refresh for intracoded blocks. In such fashion, video encoder 2 is able to 20 optimize the quality of, and reduce the number of errors in, video data transmission to video decoder 20, yet also manage the transmission bandwidth by controlling the number of intra-coded blocks that are sent to video decoder 20 over time. In some aspects, feedback 18 may include indications 25 regarding any number of different types of packet errors associated with coded video data, including indications of packet loss. For purposes of illustration only, it will be assumed in the following description that feedback 18 relates to packet loss feedback.

Video encoder 2 includes a storage medium 4, a display device 12, and one or more processors 14. Storage medium 4 may include any permanent or volatile memory that is capable of storing instructions and/or data, such as, for example, synchronous dynamic random access memory 35 (SDRAM), read-only memory (ROM), non-volatile random access memory (NVRAM), embedded dynamic random access memory (eDRAM), static random access memory (SRAM), or flash memory. Display device 12 may be any device capable of displaying video data for display pur- 40 poses, such as an LCD (liquid crystal display), plasma display device, or other television (TV) display device. Processors 14 may include one or more processors, such as a control processor, a video processor, a graphics processor, and/or a display processor, such as is shown in the example 45 of FIG. 2. Display device 12 and processors 14 are each capable of reading from and writing to storage medium 4. For example, processors 14 may execute instructions from one or more functional modules that are stored in storage medium 4. These modules may comprise downloadable 50 modules that are downloaded dynamically, over the air, into storage medium 4. In one aspect, one or more of these modules may comprise binary instructions that are generated or compiled from application programming interface (API) instructions created by an application developer.

FIG. 1 shows an example of three such modules stored within storage medium 4, each of which may comprise one or more instructions, or code, for execution. As seen in FIG. 1, storage medium 4 stores instructions for an intra-encoder 6, a predictive (inter) encoder 8, and a rate adaptation 60 module 10. Intra-encoder 6 is capable, upon execution, of intra-coding one or more segments (e.g., blocks) of video data within one or more video units (e.g., frames). Intra-coding performed by intra-encoder 6 may comprise independent coding that is not based on a coding of data 65 segments contained within other video units within a sequence of video units. Predictive encoder 8 is capable of

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inter-coding one or more segments of video data to reduce temporal redundancies across video units. Inter-coding may implement predictive techniques to encode video data based upon motion compensation from previously or subsequently coded frames. Rate adaptation module 10 is capable of receiving and processing packet loss feedback 18, which is provided by video decoder 20, to determine (e.g., set, modify) the rate of intra-coding (or the intra-refresh rate) implemented by intra-encoder 6.

Although the example of FIG. 1 shows intra-encoder 6, predictive encoder 8, and rate adaptation module 10 as modules stored within storage medium 4 and executable by processors 14, these modules may, in other cases, be implemented within hardware (such as within the hardware of processors 14). Processors 14 may, in some cases, include one or more types of processors, such as a general purpose microprocessor and/or other types of processors, such as shown in the example of FIG. 2.

Over time, video encoder 2 sends coded video data to video decoder 20 for decoding and display. If video decoder 20 determines that it has lost one or more packets of video data (such as by tracking packet sequence numbers), it may send an indication that it has lost coded video data to video encoder 2. In FIG. 1, this indication comprises packet loss feedback 18. In some cases, packet loss feedback 18 may include Real-Time Control Protocol (RTCP) data or messages. For example, packet loss feedback 18 may include one or more General Negative Acknowledgment (GNACK) messages that each may identify one or more packets, such as packets that have been lost. Packet loss feedback 18 may include one or more Picture Loss Indication (PLI) messages, which generally indicate the loss of an undefined amount of coded video data belonging to one or more pictures. In some cases, packet loss feedback 18 may include one or more Real-Time Transport Protocol (RTP) messages, such as receiver report messages, that may also provide information regarding packet loss.

In the example of FIG. 1, video decoder 20 includes a storage medium 22, a display device 28, and one or more processors 30. Similar to storage medium 4, storage medium 22 may include any permanent or volatile memory that is capable of storing instructions and/or data, such as, for example, synchronous dynamic random access memory (SDRAM), read-only memory (ROM), non-volatile random access memory (NVRAM), embedded dynamic random access memory (eDRAM), static random access memory (SRAM), or flash memory. Similar to display device 12. display device 28 may be any device capable of displaying video data for display purposes, such as an LCD (liquid crystal display), plasma display device, or other television (TV) display device. Processors 30 may include one or more processors, such as a control processor, a video processor, a graphics processor, and/or a display processor. Display device 28 and processors 30 are each capable of reading from and writing to storage medium 22.

As shown in FIG. 1, storage medium 22 may include instructions for two functional modules: an intra-decoder 24 and a predictive (inter) decoder 26. Intra-decoder 24 is capable of decoding and display intra-coded segments (e.g., blocks) of video data contained within one or more video units (e.g., frames) that are sent across network 16 by video encoder 2. Predictive decoder 26 is capable of decoding and displaying inter-coded segments that are sent by video encoder 2. Processors 30 may execute these modules 24 and 26 to display video data on display device 28. Although the example of FIG. 1 shows intra-decoder 24 and predictive decoder 26 as modules stored within storage medium 22 and

executable by processors **30**, these modules may, in other cases, be implemented within hardware (such as within the hardware of processors **30**). Processors **30** may, in some cases, include one or more types of processors, such as a general purpose microprocessor and/or other types of processors, such as shown in the example of FIG. **2**.

In some cases, video decoder 20 is further capable of performing encoding operations. In these cases, video decoder 20 may include modules similar to intra-encoder 6, predictive encoder 8, and rate adaptation module 10. In some cases, video encoder 2 is further capable of performing decoding operations, such as when receiving coded video data from device 20 (which may perform both decoding and encoding operations). In these cases, video encoder 2 may include modules similar to intra-decoder 24 and predictive decoder 26

FIG. 2 is a block diagram illustrating example processors 50 that may be included within video encoder 2 and/or video decoder 20 shown in FIG. 1, according to one aspect. When 20 included within video encoder 2, processors 50 shown in FIG. 2 may be part of processors 14. When included within video decoder 20, processors 50 may be part of processors 30. Processors 14 and processors 30 may each include one or more processors, and may, in some cases, include one or 25 more of the processors shown in the example of FIG. 2.

Processors 50 may comprise multiple programmable processors. As shown in FIG. 2, these may include a control, or general-purpose, processor 52, a video processor 54, a graphics processor 56, and a display processor 58. Control 30 processor 52 may be capable of controlling video processor 54, graphics processor 56, and/or display processor 58. In one aspect, processors 50 may include other forms of multi-media processors.

Video processor **54** may be utilized during video encoding 35 and/or decoding operations, and may implement various coding algorithms. For example, video processor **54** may implement, in some cases, H.261, H.263, Motion Picture Experts Group (MPEG) 1, MPEG2, and/or MPEG 4 video coding standards, as well as the H.264 standard and its 40 counterpart, ISO/IEC MPEG-4, Part 10, i.e., Advanced Video Coding (AVC). Video processor **54** may also implement one or more discrete cosine transform (DCT) or inverse discrete cosine transform (IDCT) operations when encoding or decoding video data.

Graphics processor **56** may be a dedicated graphics rendering device utilized to render, manipulate, and display computerized graphics. Graphics processor **56** may implement various complex graphics-related algorithms. For example, the complex algorithms may correspond to representations of two-dimensional or three-dimensional computerized graphics. Graphics processor **56** may implement a number of so-called "primitive" graphics operations, such as forming points, lines, and triangles or other polygon surfaces, to create complex, three-dimensional images on a 55 display, such as display device **12** or **28** (FIG. **1**). Display processor **58** manages the display of pixel data on the display device. Display device **58** may perform various operations on the pixel data, such as rotation and/or scaling operations.

FIGS. 3A-3F are diagrams illustrating a frame sequence, where a certain number of blocks are intra-coded within each frame of the sequence, according to one aspect. These diagrams show segments of video data as blocks (e.g., macro-blocks), but segments of any size or shape of video 65 data could be used. The blocks may comprise macro-blocks, such as 16×16 pixel macro-blocks, for example, or may, in

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some cases, comprise small sub-partitions (16×8, 8×16, 8×8, 4×8, 8×4, 4×4, etc.) pixel blocks.

In the example of FIGS. 3A-3F, it is assumed that an example frame sequence includes six frames that follow one another consecutively in a temporal sequence. (Of course, in other cases or examples, a frame sequence may include many more than six frames.) Each frame may include an array of blocks arranged at different positions. Each of FIGS. 3A-3F shows one of these six frames for purposes of illustration. Video encoder 2 (FIG. 1) is capable of coding (e.g., intra-coding and/or inter-coding) blocks of the video data within each frame and sending the frame to video decoder 20 for decoding and display. Video encoder 2 may use its predictive encoder 8 to inter-code one or more blocks of video data within any given frame, and may use its intra-encoder 6 to intra-code one or more blocks of video data within any given frame. Video encoder 2 may use its rate adaptation module 10 to adjust the rate of intra-coding (i.e., intra-refresh rate) of blocks, or the number of intracoded blocks in each frame.

In FIGS. 3A-3F, each of the temporally consecutive frames includes six intra-coded blocks 102. FIGS. 3A-3F also depict blocks that are not yet intra-coded 100, and also previously intra-coded blocks 104. Any blocks that are not intra-coded within a given frame may, in some cases, be inter-coded. However, as can be seen in the six-frame sequence of FIGS. 3A-3F, if no blocks are intra-coded twice (because each block may be unique, as is the case in this example), then all of the blocks will be intra-coded once in the six-frame sequence. It should be noted that, in this example, the same number of blocks (six) are intra-coded in each frame, but this is not necessary. The blocks that are not intra-coded in each frame may be inter-coded, and may reference one or more frames where a corresponding best-matching block is intra-coded or inter-coded.

Because, in FIGS. 3A-3F, only six blocks are intra-coded in each frame of the sequence, it takes six frames to complete the intra-coding of each block at all positions in the entire frame. However, there may be interdependency between frames, given that other blocks are inter-coded, meaning that coded bit streams may be susceptible to propagating channel errors, which may result in errors within the video frames decoded by video decoder 20. Increasing the number of intra-coded blocks in each transmitted frame may reduce errors and error propagation caused by channel errors. However, increasing the number of intra-coded blocks in each frame may also increase transmission bandwidth, given that more bits may be used for intra-coding than for inter-coding.

Thus, rate adaptation module 10 of video encoder 2 may dynamically determine, or adjust, the rate of intra-coding that is performed by intra-encoder 6. For example, rate adaptation module 10 may determine the number of intracoded blocks for each frame in the sequence, and thus may increase or decrease this number in an attempt to balance the bandwidth and error reduction parameters. As previously described, rate adaptation module 10 may determine the intra-refresh rate based, in some cases, upon packet loss feedback 18 that is received from video decoder 20. For example, if video decoder 20 indicates that there is a significant amount, or increase, in packet losses or errors, rate adaptation module 10 may determine to increase the intra-refresh rate to help reduce the possibility of further propagating errors. This will be described in more detail below.

FIG. 4 is a flow diagram illustrating a method that may be performed by a video encoder, such as video encoder 2

shown in FIG. 1, to determine a rate for intra-coding segments of video data based upon received error feedback, such as error feedback 18, according to one aspect. FIG. 4 shows the details of a high-level method that may be implemented by a video encoder, and further details of 5 various aspects are also shown in subsequent figures, such as FIGS. 5-8. In the following description of FIG. 4, it will be assumed that the method is performed by video encoder 2 for purposes of illustration only.

Video encoder 2 may receive feedback 18 from video 10 decoder 20 that indicates an error (e.g., loss of data) associated with coded video data, such as coded video data that may have been previously sent by video encoder 2 to video decoder 20 (106). Upon receipt of feedback 18, video encoder 2 may determine a rate for intra-coding a specified 15 number of segments (e.g., blocks) of video data for each video unit within a sequence of video units (e.g., frames), wherein an intra-coding of segments of video data within one video unit is independent of a coding of segments of video data within other video units within the sequence 20 (108). Video encoder 2 may repeat the acts of processing received feedback 18 and determining a rate for intra-coding for multiple video units and/or sequences of video units.

In some cases, feedback 18 that is sent from video decoder 20 indicates a loss of at least a portion of the coded 25 video data by video decoder 20. Each segment of video data may comprise one or more blocks of video data, and each video unit within the sequence of video units may comprise one or more frames or slices of video data. An intra-coding of blocks of video data within one frame may be independent of a coding of blocks of video data within other frames within the sequence.

Feedback 18 may, in one aspect, comprise at least one of a picture loss indication (PLI) and a general negative acknowledgment (GNACK). Feedback 18 may comprise 35 feedback indicating a number of lost packets containing coded video data, each lost packet having been previously sent to, but not received by, video decoder 20 by video encoder 2. In some cases, determining the rate for intracoding may be based, at least in part, on a target correction 40 time for correcting errors across one or more video units within the sequence, as will be described in further detail below.

When feedback 18 includes a GNACK that indicates a number of lost packets, determining the rate for intra-coding 45 may be based, at least in part, on the number of lost packets. In some cases, video encoder 2 may also receive a report from video decoder 20 indicating a second number of lost packets. For example, this report may be one of a series of RTCP Receiver Reports periodically sent by video decoder 50 20 to video encoder 2. As will be described in further detail below, video encoder 2 may, in one aspect, calculate a packet error rate based upon at least one of the number of lost packets indicated by the GNACK and the second number of lost packets indicated by the report, wherein determining the 55 rate for intra-coding is based, at least in part, on the packet error rate. For example, video encoder 2 may calculate a packet error rate based upon which of the GNACK or report indicates a higher volume of lost packets. Determining the rate of intra-coding may be further based, at least in part, on 60 an estimated number of packets contained within each video unit within the sequence of video units.

If error feedback 18 includes a GNACK message, the GNACK may identify each lost packet, and video encoder 2 may calculate an elapsed time based upon a difference 65 between a current time and a time at which one of the identified lost packets was previously sent to video decoder

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20, as is described in further detail below in reference to FIG. 8. Video encoder 2 may determine the rate of intracoding based, at least in part, on this elapsed time. In some cases, each lost packet identified by the GNACK has an associated sequence number. Sequence numbers may be initially assigned to packets, in ascending order, when the packets are sent from video encoder 2 to video decoder 20. Video decoder 20 may keep track of which packets it has received, and also the corresponding sequence numbers of these packets. Thus, the GNACK sent by video decoder 20 may identify the sequence numbers of any lost packets (or sequence number ranges for lost packets), assuming the sequence numbers are in ascending order.

Upon receipt of a GNACK that identifies lost packets by sequence number, video encoder 2 may calculate the elapsed time. Video encoder 2 may determine which lost packet identified by the GNACK has a sequence number that is greater than a sequence number of a first packet of a prior sequence of video units having intra-coded segments that was previously sent by video encoder 2 to video decoder 20. Video encoder 2 may then calculate the elapsed time based upon a difference between the current time, as maintained by video encoder 2, and a time at which the determined lost packet was previously sent to video decoder 20 from video encoder 2, as is described in more detail with reference to FIG. 8. Video encoder 2 may make such a calculation of the elapsed time by keeping track of when each packet, as identified by sequence number, was sent to video decoder 20.

As described above, video encoder 2 may repeat the receiving of feedback 18 and the determining of the rate for intra-coding for each video unit in multiple sequences of video units. In some cases, video encoder 2 may determine whether to repeat the processing of received feedback 18 and the determining of the rate for intra-coding in multiple sequences of video units based upon whether a detected amount of motion across different video units within one sequence of video units exceeds a motion threshold, as is described in further detail below. In one aspect, the motion threshold is an adjustable parameter that may, in some cases, be set to a higher threshold value by video encoder 2 when video encoder 2 uses a higher rate of intra-coding segments of video data.

FIG. 5 is a flow diagram illustrating a method for determining an intra-refresh rate for intra-coding blocks of video data based upon received packet loss information, according to one aspect. Although the example of FIG. 5 refers to blocks, or macro-blocks, of video data, various other segments (sizes, shapes) of video data may be intra-coded in various other scenarios. In one aspect, the method shown in FIG. 5 may be performed by video encoder 2 (FIG. 1) based upon packet loss feedback 18 that is received from video decoder 20.

In various scenarios, video encoder 2 (and, more particularly, rate adaptation module 10) may determine a macroblock intra-refresh rate based upon received packet loss feedback, and may use various criteria or factors to calculate the intra-refresh rate. For example, as will be described in further detail, video encoder 2 may use a target correction time, a probability of uncorrected errors, a number of errors detected, an amount of motion, and/or a limited duration of user-perceived error when calculating an intra-refresh rate.

Various parameters are described with reference to FIG. 5 and subsequent flow diagrams. Some of these parameters are defined below:

Target\_Correction\_Time (in seconds): maximum amount of time after detection of an error by video encoder 2, such as a packet loss, for which the video encoder 2 has to correct the error.

Current\_Frame\_Rate (in frames per second): the current 5 number of frames per second used by video encoder 2. If the video encoder 2 varies this parameter with changes in the target\_enc\_rate (target encode rate) parameter, then the rate adaptation module 10 may estimate this from delivery of frames from the video encoder 2. In some cases, the video 10 encoder 2 may indicate the value of Current\_Frame\_Rate to the rate adaptation module 10.

Max\_Intra\_MB\_% (0-100%): maximum allowable percentage of intra-coded macro-blocks. This parameter may be used to limit degradation in video quality since the frame 15 size may be fixed according to data and frame rate.

Intra\_Repeat (integer): number of times a full macroblock intra-refreshed sequence (across a sequence of frames) is to be performed across the entire image when a packet loss is detected.

Intra\_MB\_%\_No\_Loss (0-100%): percentage of intracoded (refreshed) macro-blocks when there is no detected packet loss.

Intra\_MB\_% (0-100%): the percentage of intra-coded (refreshed) macro-blocks that the video encoder **2** is to insert 25 per frame (or slice).

Num\_Intra\_Frames (integer): the number of frames that the video encoder **2** has intra-coded (refreshed) macroblocks.

In addition, various formulas and calculations that are 30 shown below include Max and Min operations. It is assumed that a Max operation selects the maximum value of two or more possible values within a set, and that a Min operation selects the minimum value of two or more possible values within a set.

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The method shown in FIG. 5 repeatedly checks to determine whether video encoder 2 has received RTCP feedback from video decoder 20 (110). Upon receipt of such feedback, video encoder determines whether the feedback indicates a loss of one or more packets that were transmitted by video 40 encoder 2 to video decoder 20 (112). For example, packet loss feedback 18 (FIG. 1) sent by video decoder 20 to video encoder 2 may indicate such a loss of one or more packets of coded video data sent by video encoder.

If video encoder **2** has not received any packet loss 45 indication, it may use rate adaptation module **10** to determine, or set, an intra-refresh rate, or percentage, for the macro-blocks (Intra\_MB\_%) equal to Intra\_MB\_%\_No\_Loss, which may comprise an adjustable parameter (**114**). In some cases, video encoder **2** may 50 automatically set or adjust Intra\_MB\_%\_No\_Loss. In some cases, a user may set or adjust this parameter, such as to increase or decrease the number of intra-coded macro-blocks in each frame.

If, however, video encoder **2** has received a packet loss 55 indication, such as packet loss feedback **18**, it may determine whether it has received a GNACK or a PLI (**116**). A GNACK may indicate one or more specific packets that were lost, or that had errors. The Intra\_refresh routine begins upon receipt of a packet loss indication (**118**).

Initially, rate adaptation module 10 of video encoder 2 may set (120) Intra\_MB\_% equal to

Min{Max\_Intra\_MB\_%,100%/(Target\_Correction\_ TimexCurrent\_Frame\_Rate)}.

Max\_Intra\_MB\_% and Target\_Correction\_Time are each adjustable parameters that may have initial, predefined val-

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ues that are optimized for the particular application or operations that are performed by video encoder 2. These parameters may also have different values for different applications (e.g., video telephony, video streaming, gaming) that are executed by video encoder 2. Rate adaptation module 10 uses the values of each of these parameters when determining the value of Intra\_MB\_%.

In some cases, rate adaptation module 10 may implement one or more optional enhancement functions, such as are shown in FIGS. 6-8, to obtain a modified value of Intra\_MB\_% (122), as will be described in more detail below. Rate adaption module 10 requests video encoder 2 to start an intra-refresh procedure for a sequence of video frames using Intra\_MB\_% as the intra-refresh rate for macro-blocks in each frame, such as by executing intraencoder 6 and predictive encoder 8 (124). Video encoder 2 also may initialize the value of Num\_Intra\_frames to zero. Input is then provided by video encoder 2 that the intrarefresh procedure has started. Video encoder 2 also indicates when each video frame boundary has been reached (126). This may be part of the same indication that is used to set a timestamp and/or market bit in RTP. When a video frame boundary has been reached, rate adaptation module 10 increments Num\_Intra\_Frames (128) which is a parameter that keeps track of the number of video frames that include intra-coded macro-blocks using the current value of Intra\_MB\_%.

At each frame boundary, rate adaptation module  ${\bf 10}$  determines  $({\bf 130})$  whether

Num\_Intra\_Frames>=Ceil[100%/Intra\_MB\_%].

If Num\_Intra\_Frames is less than this ceiling (or maximum threshold), then video encoder 2 checks to determine whether it has received another packet loss indication from video decoder 20 (131). If not, video encoder 2 again checks for a video frame boundary (126). If, however, video encoder 2 has received another packet loss indication from video decoder 20, such as packet loss feedback 18, video encoder 2 may determine whether it has received a GNACK or a PLI (116). In this case, video encoder 2 then continues with the Intra\_refresh routine (118).

If Num\_Intra\_Frames is greater than or equal to this ceiling, video encoder 2 has sent a full sequence of video frames to video decoder 20 using the current value of Intra\_MB\_% for each frame within this sequence (whereby an indicated number of macro-blocks in each frame have been intra-coded). Rate adaptation module 10 then performs the Repeat\_Intra routine shown in FIG. 5 (132). In this routine, rate adaptation module 10 determines whether the number of times that video encoder 2 has intra-refreshed an entire sequence of video frames (Intra-refreshed) is greater than, or equal to, the value of Intra\_Repeat (which is another adjustable parameter that may have an initial value of two) (134). The value of Intra\_Repeat may be adjusted to account for possible motion associated with one or more macroblocks within an intra-refreshed frame sequence. Such motion may potentially result in errors that can be corrected by repeating the intra-refresh sequence for a video sequence 60 one or more times.

Rate adaptation module 10 returns to a no-loss-detected operation, and sets Intra\_MB\_%=Intra\_MB\_%\_No\_Loss, when the number of times that a complete intra-refresh sequence has been performed (Intra-refreshed) is greater than, or equal to, the value of Intra\_Repeat. If, however, Intra-refreshed is less than the value of Intra\_Repeat, then rate adaptation module 10 increments the value of Intra-

refreshed (136) and again performs the Intra\_refresh routine (118) for a subsequent sequence of video frames, as is shown in FIG. 5.

FIG. 6 is a flow diagram of a method that may be performed by video encoder 2, shown in FIG. 1, to obtain a modified value of Intra\_MB\_%. This modified value of Intra\_MB\_% may be used within the method of FIG. 5, as specified within the box "Optional enhancement of Intra\_MB\_%" (122) in FIG. 5. In one aspect, video encoder 2 may perform the method shown in FIG. 5 to calculate an initial value of Intra\_MB\_%, but may then additionally perform the method shown in FIG. 6 to determine a modified value of Intra\_MB\_% for use when intra-coded macroblocks within each video frame of a frame sequence.

By performing the method shown in FIG. **6**, video encoder **2** may limit the probability of having uncorrected errors after one intra-refresh sequence of intra-coded macroblocks within the frame sequence. Rate adaptation module **10** may determine a modified value of Intra\_MB\_% based, 20 at least in part, on an estimate of an RTP packet loss rate during the intra-refresh sequence, and may make use of the fraction\_lost parameter that is included within an RTCP Receiver Report that is provided by video decoder **20**, indicating the fraction of packets that have been lost by 25 video decoder **20** since the last Receiver Report was sent.

In addition, rate adaptation module 10 may be able to increase the value of Intra\_MB\_% when it has detected a larger number of errors (e.g., packet losses) from video decoder 20. By doing so, video encoder 2 may be able to more quickly correct errors that may be seen by the user of video decoder 20 (as displayed on display device 28).

Various parameters are described with reference to FIG. 6 and subsequent flow diagrams. Some of these parameters are defined below:

Target\_Err\_Pr (0-1): target probability that there will still be an error in the image after a single intra-refresh sequence (i.e., sequence of video frames). In the example of FIG. **6**, this parameter is set higher than an expected PER (packet 40 error rate) times the number of packets per frame.

 $\alpha$  (0-1): IIR (infinite impulse response) filter update factor for estimating avg\_norm\_ppf (average normalized packets per frame). Video encoder 2 may start with a value of 0.1.  $\beta$  (0-10): tuning factor for determining Intra\_MB\_%. 45

Video encoder 2 may start with a value of 1.0.

PER (0-1): packet error rate, which may be set equal to the fraction\_lost parameter from the most recent RTCP Receiver Report received from video decoder 20.

avg\_norm\_ppf (real number): average 'normalized' pack- 50 ets per frame, which may be normalized to a target encode rate of 50 kbps (kilobits per second) and 10 fps (frames per second).

inst\_ppf (integer): packets per frame in current frame
est\_ppf (real number): estimated number of packets per 55
future frame

Target\_enc\_rate (bits per second): current value of target\_encode\_rate set by rate adaptation module 10.

As is shown in FIG. 6, video encoder 2 may calculate, or estimate, the value of avg\_norm\_ppf (150) when calculating 60 a modified value of Intra\_MB\_% (163). Once a packet has been prepared and is ready for transmission by video encoder 2 (152), video encoder 2 may increment inst\_ppf, which is number of packets per frame in the current frame (154). In some cases, inst\_ppf indicates the number of RTP 65 packets in the frame inst\_ppf is further incremented for each packet within the frame until a final packet indicates the end

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of the frame (156). At this point, video encoder 2 is able to calculate (158) an updated value of avg\_norm\_ppf as follows:

avg\_norm\_ppf=(1-\alpha)\timesavg\_norm\_ppf+(\alpha\timesinst\_ppf)\times (50 kbps/target\_encode\_rate)\times(Current\_Frame Rate/10 fos).

wherein it may be assumed, in this particular example, that the target encode rate is normalized with respect to a predefined value of 50 kbps, and that the current frame rate is normalized with respect to a predefined value of 10 fps. These predefined values are used in certain example scenarios only, and may have various different values in different configurations. The value of inst\_ppf can then be reset to zero for the next frame (160). As shown, the value of avg\_norm\_ppf is based upon the values of inst\_ppf, target\_encode\_rate, and Current\_Frame\_Rate (which is the frame rate of the current frame). As shown in FIG. 6, this value of avg\_norm\_ppf may then be used to calculate a value of est\_ppf.

In one aspect, the value of est\_ppf (estimated packets per frame) may be calculated as follows:

est\_ppf=avg\_norm\_ppfx(target\_enc\_rate/50 kbps)x (10 fps/Current\_Frame\_Rate),

where, in this example, the target encode rate is normalized with respect to a predefined value of 50 kbps, and that the current frame rate is normalized with respect to a predefined value of 10 fps (166). The predefined values may be optimized, and may have different values in various different configurations. The target encode rate (target\_enc\_rate) may be provided as one input parameter to the equation, and the value of Intra\_MB\_% from the base algorithm shown in FIG. 5 may be provided as another (164). In one aspect, the target encode rate, target\_enc\_rate, may comprise an adjustable parameter that may be tuned based upon, for example, the type of application that is being executed by video encoder 2. In one aspect, the target encode rate may be tuned based upon a dynamic rate adaptation algorithm, where video encoder 2 adapts its rate to network conditions.

Rate adaptation module 10 of video encoder 2 may set the value of PER (packet error rate) equal to the value of fraction\_lost that is contained within an RTCP Receiver Report sent by video decoder 20, which, as described above, indicates the fraction of packets that have been lost by video decoder 20 since the last Receiver Report was sent (168). Rate adaptation module 10 may then calculate an updated value of Target\_Err\_Pr, which is the target probability that there will still be an error in the image after a single intra-refresh sequence, as follows (170):

Target\_Err\_Pr=Max{Target\_Err\_Pr, est\_ppf\*PER}.

After making this calculation, rate adaptation module 10 is capable of calculating an updated, or modified, value of Intra\_MB\_% based upon an initial value of Intra\_MB\_% as determined by the method, or base algorithm, shown in FIG. 5. This initial value of Intra\_MB\_% as determined by the method of FIG. 5 is indicated as Intra\_MB\_% from Base algo (such as in FIG. 6, where the base algorithm is the one shown in FIG. 5). The updated value of Intra\_MB\_% may be calculated (172) as follows:

Min{Max\_Intra\_MB\_%,Max{βx100%xlog(1-PER)x est\_ppf/log(1-Target\_Err\_pr), Intra\_MB\_% from Base Algorithm}.

This updated value of Intra\_MB\_% may then be used by the base algorithm shown in FIG. 5, when rate adaptation

module 10 intra-codes macro-blocks within frames of a sequence according to this updated value of Intra\_MB\_% (174).

The method shown in FIG. 6 calculates an updated value of Intra\_MB\_% to help video encoder 2 limit the probability of having uncorrected errors after one intra-refresh sequence of intra-coded macro-blocks within the frame sequence. As a result, the updated, or modified, value of Intra\_MB\_% will typically be an increased value, such that intra-refreshing of macro-blocks occurs more quickly, and in fewer frames within a sequence. In obtaining the values and using the calculations shown in FIG. 6, it may, in some cases, be assumed that every slice of a frame (where one slice may contain one or more macro-blocks) contains at least one intra-coded macro-block.

Using this assumption, it may further be assumed that any packet loss during an entire intra-refreshed frame sequence (have intra-coded macro-blocks refreshed at a rate of Intra\_MB\_%) may cause errors in some macro-blocks because certain intra-coded macro-blocks may be lost. 20 These assumptions may comprise a conservative assumption about how quickly to intra-refresh a frame sequence. Assuming that the probability of packet errors is distributed throughout the frame sequence, all packets throughout the sequence may have an equal probability of causing uncorrected errors. Thus, the probability of there being no error after one intra-refreshed frame sequence may be equal to  $(1-PER)^P$ , where P is the number of packets in the entire frame sequence. Therefore, the total number of packets in an intra-refresh sequence may be limited to less than

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log(1{-}Target\_Err\_Pr)/log(1{-}PER).
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The number of packets in an intra-refreshed (IR) frame sequence may then be calculated as

# frames in IR sequencexavg # packets per frame during IR sequence, where # frames in IR sequence=100%/Intra\_MB\_%.

As a result, the estimated average number of packets per frame during an IR sequence may be calculated as

est\_ppf=Avg\_norm\_ppfx(target\_enc\_rate/50 kbps)x (10 fps/Current\_Frame\_Rate).

Combining these results provides the following formula with respect to the modified value of Intra\_MB\_% that is determined by the method shown in FIG. 6:

$$\begin{split} \log(1-\text{Target\_Err\_Pr}) / \log(1-\text{PER}) > = & (100\%/\text{In-tra\_MB\_\%}) \times \text{est\_ppf}. \end{split}$$

FIG. 7 is a flow diagram of a method that may be performed by rate adaptation module 10 of video encoder 2 50 (FIG. 1) to calculate a packet error rate (PER) upon receipt of a GNACK message, according to one aspect. As described in reference to the method of FIG. 6 to calculate a modified value of Intra\_MB\_%, a value of PER was selected and set equal to the value of the fraction\_lost 55 parameter included within a Receiver Report sent by video decoder 20. FIG. 7 shows a method that may be used by rate adaptation module 10 to obtain a modified value of PER based upon the receipt of a GNACK.

As shown in FIG. 7, rate adaptation module 10 is capable 60 of calculating a value of a parameter GNACK\_PER based on the number of packets indicated in a received GNACK message (180). The GNACK is sent from video decoder 20. Video decoder 20 may determine the number of packets within one or more frames have errors (e.g., have been lost), 65 and may indicate this number within the GNACK. In addition, within the GNACK, video decoder 20 may identify

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the specific packets, by sequence number, that have errors. In one aspect, these packet sequence numbers are initially determined by video encoder 2 when the packets are first sent to video decoder 20 via network 16.

Upon receipt of the GNACK, rate adaptation module 10 may compare the value of GNACK\_PER with the value of fraction\_lost that was included within the prior RTCP Receiver Report sent from video decoder 20 (182). If GNACK\_PER is greater than, or equal to, fraction\_lost, then rate adaptation module 10 sets the value of PER equal to GNACK\_PER (184). If, however, GNACK\_PER is less than fraction\_lost, then rate adaptation module 10 sets the value of PER equal to the value of fraction\_lost from the last Receiver Report (186).

Thus, as can be seen from the method of FIG. 7, rate adaptation module 10 selects a value of the packet error rate, PER, as the greater of GNACK\_PER and fraction\_lost. By selecting the greater of these two values, rate adaptation module 10 uses a higher value of PER when calculating the modified intra-refresh rate, Intra-MB\_%, within the method of FIG. 6. In this aspect, rate adaptation module 10 increases Intra-MB\_% when a higher volume of packet errors are indicated, such as by the fraction\_lost parameter or by a GNACK message.

FIG. 8 is a flow diagram illustrating a method that may be performed by rate adaptation module 10 of video encoder 2 (FIG. 1) to determine an updated intra-refresh rate based upon a receipt of one or more GNACK messages from video decoder 20 that may each include packet sequence numbers,
according to one aspect. By performing the method of FIG. 8, video encoder 2 may try and limit the total time that a user of video decoder 20 may perceive an error. Rate adaptation module 10 of video encoder 2 may use larger values of Intra\_MB\_% for older errors, such that these older errors
may be corrected more quickly.

Rate adaptation module 10 may maintain a mapping of RTP packet sequence numbers to points in time, or time-stamps, at which point the packets were sent from video encoder 2 (190). For example, video encoder 2 may send a first packet having a first sequence number at a first point in time (e.g., 100 ms), and may send a second packet having a second sequence number at a second point in time (e.g., 110 ms). Rate adaptation module 10 may keep a mapping of the sequence numbers-to-timestamps for each of these two packets, as well as any other packets that are sent from video encoder 2 to video decoder 20.

When video encoder 2 receives a GNACK from video decoder 20, indicating that one or more packets have an error (e.g., have been lost), video encoder 2 is capable of determining the sequence numbers of each of the packets identified within the GNACK. For example, if the GNACK indicates that four different packets have errors, video encoder 2 is capable of determining the RTP sequence numbers of each of these four packets. These sequence numbers may be associated with these four packets, and included within the GNACK. (As described previously, these RTP sequence numbers are the sequence numbers originally assigned to the packets when they were first sent by video encoder 2 to video decoder 20. In one aspect, the sequence numbers are assigned by video encoder 2 in ascending order for packets that are sequentially sent to video decoder 20.)

Upon determination of the sequence numbers of the packets identified by the received GNACK, video encoder 2 may the determine (192) the earliest, or smallest, sequence number for an identified packet (RTP\_SN) that is greater than the sequence number of the first packet of the last

intra-refreshed frame sequence containing intra-coded macro-blocks (RTP\_SN\_Last\_IR\_Sequence), or

RTP\_SN>RTP\_SN\_Last\_IR\_Sequence.

Video encoder **2** uses its mapping information to determine the timestamp for the packet identified within the GNACK having this earliest sequence number of RTP\_SN. This timestamp corresponds to the point in time at which video encoder **2** originally sent this packet having sequence number RTP\_SN to video decoder **20**. This timestamp is therefore based upon the clock maintained by video encoder **2**. Video encoder **2** compares this timestamp (TS\_of\_earliest\_GNACK\_RTP\_SN) to the current time (current\_time), as indicated by the clock of video encoder **2**, and calculates (**194**) a value of the difference (Time\_elapsed) as follows: 15

Time\_elapsed=Current\_Time-TS\_of\_earliest\_G-NACK\_RTP\_SN.

The value of the Time\_elapsed parameter indicates the period of time that has elapsed since the packet having sequence number RTP SN has been sent by video encoder 2. Rate adaptation module 10 can then modify the value of Intra\_MB\_% (used for intra-coding macro-blocks in subsequent intra-refreshed sequences) based upon this elapsed time period. Because RTP\_SN is the earliest, or smallest, sequence number for an identified packet in the GNACK that is greater than the sequence number of the first packet of the last intra-refreshed frame sequence containing intracoded macro-blocks, rate adaption module 10 attempts to modify the value of Intra MB % to try and correct older errors (which occurred longer ago) more quickly. Typically, rate adaption module 10 will increase the value of Intra\_MB\_% in these cases, and then use this modified value within the method shown in FIG. 5 (in the Intra\_refresh

In one aspect, rate adaptation module 10 may calculate  $^{35}$  (196) the modified value of Intra\_MB\_% as follows:

Set Intra\_MB\_%=Min{Max\_Intra\_MB\_%,100%/ [(Target\_Correction\_Time-Time\_elapsed)xCurrent Frame\_Rate]}.

As with certain of the formulas or calculations described previously, rate adaptation module 10 may bound the value of Intra\_MB\_% by Max\_Intra\_MB\_% in certain cases. The value of Max\_Intra\_MB\_% may be configured or adjusted by video encoder 2 or by a user, based upon the type of 45 application executed by, or environment provided, by video encoder 2. In the case above, rate adaptation module 10 also may determine the modified value of Intra\_MB\_% based, at least in part, upon a difference between Target\_Correction\_ Time and Time\_elapsed. The Target\_Correction\_Time 50 parameter, as described previously, indicates the maximum amount of time after detection of an error by video encoder 2, such as a packet loss, for which the video encoder 2 has to correct the error. In this case, this maximum correction time becomes a targeted maximum amount of time that a 55 receiver sees a propagation of the error. The value of Target\_Correction\_Time may be configured or changed for purposed of optimization.

In some aspects, video encoder 2 may implement another enhancement for determining a modified value of 60 Intra\_MB\_% that may be used within the Intra\_refresh routine shown in FIG. 5. In these aspects, rate adaptation module 10 may attempt to measure an amount of motion over a past video frame sequence during which Intra\_MB\_% was used upon detection of one or more packet errors. As 65 described previously in reference to FIG. 5, rate adaptation module 10 may potentially repeat an intra-refresh sequence

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for a frame sequence one or more times, according to the value of an adjustable parameter Intra\_repeat. The value of this parameter may be adjusted based upon detected motion, which may be a function of motion vectors over the frames within a frame sequence.

Rate adaptation module 10 may compare the amount of detected motion to a motion threshold value, which may comprise a value that is configured or adjusted by a user or by video encoder 2. For example, if a higher value of Intra\_MB\_% is being currently used, video encoder 2 may use a higher threshold value, given that intra-refreshing is occurring at a higher rate. The threshold may apply to motion statistics for an overall frame, or sequence, but may also apply to motion for specific macro-blocks, as well.

If the amount of detected motion is less than this threshold, then rate adaptation module 10 may modify the value of Intra\_MB\_% to be set equal to the no-loss value (as used in FIG. 5), according to:

Intra\_MB\_%=Intra\_MB\_%\_No\_Loss.

However, if the amount of motion exceeds this threshold, rate adaptation module 10 may repeat the intra-refresh procedure of an entire video frame sequence, and/or may increase the value of Intra\_repeat, such that an additional refresh can be performed. The repeating of the intra-refresh sequence may help correct any errors that may be caused by excessive motion within one or more frames.

FIG. 9 is a flow diagram illustrating a method that may be performed by video decoder 20 (FIG. 1) to determine whether to send a PLI (picture loss indication) message or another indication, such as a General Negative Acknowledgment (GNACK) to video encoder 2, according to one aspect. As described previously, video decoder 20 is capable of sending error (e.g., packet loss) feedback 18 to video encoder 2 upon detection of one or more packet errors by intra-decoder 24 and/or predictive decoder 26. This feedback 18 may include one or more GNACK and/or PLI messages. GNACK messages may specifically identify packets that have errors, while PLI messages may more generally identify decoding errors without necessarily identifying packets.

However, there may be certain disadvantages to sending too many PLI messages, for example, to video encoder 2 over time. For example, video decoder 20 may send a first PLI to video encoder 2, which may trigger rate adaptation module 10 of video encoder 2 to begin an intra-refresh sequence, or to increase the rate of intra-encoding macroblocks within a frame sequence by increasing the value of Intra\_MB\_%. If the video decoder 20, shortly thereafter, sends another PLI to video encoder 2 before receiving any of the newly refreshed macro-blocks, video encoder 2 may again, but potentially unnecessarily, increase the value of Intra\_MB\_%, assuming that there are new, or additional, errors perceived by video decoder 20. In these types of scenarios, PLI messages from video decoder 20 may cross paths with intra-refreshed macro-blocks that are sent by video encoder 2 to correct previously indicated errors. Thus, in some cases, it may be more beneficial for video decoder 20 to perform certain checks before sending repeated PLI messages to video encoder 2 over a short period of time.

According to the method of FIG. 9, video decoder 20 may wait for decoder output from intra-decoder 24 and/or predictive decoder 26 (200). Video decoder 20 may then determine whether any decoding errors were detected by intra-decoder 24 and/or predictive decoder 26 (202). In some cases, video decoder 20 may detect certain errors that are not detected by the RTP layer, such as through analysis of the bit

stream of the coded video data. If a decoding error is detected, such as an error in one of the received, coded packets, or even the loss of one or more packets, video decoder 20 may determine the current time (current\_time) according to its clock.

Video decoder 20 may also determine, or lookup, a value for RTT, which is the value of the round-trip travel time for data communication between video decoder 20 and video encoder 2. In some cases, the value of RTT may be initially set, or configured, within video decoder 20 as a default 10 value. Video decoder 20, however, is capable of dynamically adjusting, or updating, the value of RTT over time based upon real-time calculations for data communication between video decoder 20 and video encoder 2. For example, video decoder 20 may be able to use information contained within 15 one or more RTCP Receiver and/or Sender Reports in order to measure, or estimate, the value of RTT. In one aspect, video decoder 20 may measure or estimate the value of RTT based upon at least one pair of messages that are exchanged between the video decoder 20 and video encoder 2.

In the example of FIG. 9, video decoder 20 may check to determine (204) whether the following condition is satisfied:

current\_time>RTT+Time\_Last\_Sent\_PLI,

where is Time\_Last\_Sent\_PLI equal to the time (according 25 to the clock of video decoder 20) that video decoder 20 last sent a PLI to video encoder 2. If the current time is, in fact, greater than the RTT plus Time\_Last\_Sent\_PLI, then video decoder 20 may send another PLI message to video encoder 2, and also set the value of Time\_Last\_Sent\_PLI to the 30 current time, according to the clock of video decoder 20 (206).

If, however, the current time is less than, or equal to, RTT plus Time Last Sent PLI, then video decoder 20 will suppress sending another PLI message to video encoder 2. In 35 these scenarios, video decoder 20 determines that it has previously sent a PLI within one round-trip travel time, and does not need to send another PLI. Video decoder 20 assumes, in these scenarios, that video encoder 2 will begin a new intra-refresh sequence, or may even adjust the value 40 of Intra\_MB\_%, in response to receiving a prior PLI sent by video decoder 20, and that video decoder 20 will, at the very earliest, begin receiving packets (and macro-blocks) within this refreshed sequence within one round-trip travel time after having sent the prior PLI. Thus, if video decoder 20 45 identifies any errors within this one round-trip travel time window, it assumes that such errors may be corrected by the intra-refresh sequence that is initiated by video encoder 2 upon receipt of the prior PLI, and that it does not need to send another PLI within this time window. Video decoder 20 50 may continue to perform the method shown in FIG. 9 for all coded packets received from video encoder 2.

FIG. 10 is a flow diagram of a method that may be performed by video decoder 20 (FIG. 1) to determine whether to send a GNACK or a PLI to video encoder 2, 55 according to one aspect. In some cases, video decoder 20 may determine to send a general PLI message to video encoder 2 if, for example, there are a large number of packets that have errors or have been lost since last sending a PLI. In one aspect, video decoder 20 may perform the 60 method shown in FIG. 10 if no prior GNACK or PLI, for example, has been sent within one round-trip travel time, such as shown in FIG. 9.

Thus, as is shown in FIG. 10, video decoder 20 may wait for decoding output (210), and determine whether processing has reached a frame boundary (or whether a new frame has been detected) (211). If so, video decoder 20 may reset

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the value of a variable called Number\_of\_packets\_lost to zero. As described below, video decoder 20 uses the value of this variable in determining whether to send a GNACK or a PLI upon detection of packet errors. When a new frame has been detected, video decoder 20 may reset the value of this variable to zero.

Video decoder 20 may identify that a particular packet sent by video encoder 2 has been lost, and may then increment the Number\_of\_packets\_lost variable (212). Video decoder 20 may then compare the value of Number\_of\_packets\_lost to a packet loss threshold value. This threshold value may comprise an adjustable parameter that may be configured, or tuned, by a user or by video decoder 20. In some instances, the value of this threshold value may depend on the type of device (e.g., type of device for video decoder 20 and/or video encoder 2), and/or the type of application that is executed (e.g., video telephony application, game application).

In one aspect, video decoder 20 may determine (214) whether the following condition is satisfied:

Number\_of\_packets\_lost>=%\_Packet\_Loss\_Threshold)xavg\_ppf, where avg\_ppf is equal to the average number of packets per frame received by video decoder 20. Video decoder 20 is capable of estimating the value of avg\_ppf by continuously determining the number of packets in each frame received by video encoder 2. If there have been a high number of packets lost since last sending a PLI, such that the number of packets lost exceeds the threshold value (which may, as shown above, also be based on avg\_ppf), video decoder 20 may send another PLI message to video encoder 2 to initiate a full intra-refresh sequence (218). Video decoder 20 may also reset the value of Number\_of\_packets\_lost to zero. In one aspect, prior to sending the PLI, video decoder 20 may also perform the method shown in FIG. 9 to determine whether a prior PLI has already been sent within one round-trip time window, and may refrain from sending another PLI when such is the case, as was previously

If, however, video decoder 20 determines that the number of packets lost since transmission of the last PLI does not exceed the threshold, video decoder 20 may send a GNACK, rather than a PLI, to video encoder 2 (216). The GNACK may identify specific packets that have errors (or that were lost). Video encoder 2 may then, in some cases, use the sequence numbers associated with the packets identified in the GNACK to perform an enhanced method of updating the intra-refresh rate, such as the method shown in FIG. 8.

In view of the description above and the methods shown in FIGS. 9-10, a video decoder, such as video decoder 20, may, in one aspect, detect at least one error associated with coded video data sent from a video encoder, such as video encoder 2. Video decoder 20 may determine a round-trip travel time for data communication between video decoder 20 and video encoder 2, determine a last-sent time at which video decoder 20 last sent video encoder 2 a feedback message indicating a previous error that was detected by video decoder 20, and compare a current time to a sum of the last-sent time and the round-trip travel time. Based upon the comparing, video decoder 20 may determine whether to send video encoder 2 a second feedback message indicating the at least one error detected by video decoder 20.

In some cases, video decoder 20 may send the second feedback message if the current time exceeds the sum of the last-sent time and the round-trip travel time, and may set the last-sent time equal to the current time. Determining the round-trip travel time may include measuring or estimating

the round-trip travel time based upon at least one message that is exchanged between video decoder 20 and video encoder 2. In some cases, video decoder 20 may receive the coded video data from video encoder 2, and may detect the at least one error at least by detecting at least one decoding of the coded video data by the video decoder.

In one aspect, video decoder 20 may detect the at least one error at least by detecting a packet loss. Upon detecting the packet loss, video decoder 20 may increment a count of lost 10 packets, compare the count of lost packets to a threshold, and send the second feedback message to video encoder 2 if the count of lost packets exceeds the threshold and if the current time exceeds the sum of the last-sent time and the round-trip travel time. If the count of lost packets does not 15 exceed the threshold, video decoder 20 may send another type of feedback message indicating the packet loss to video encoder 2.

In some cases, upon sending the second feedback message to video encoder 2, video decoder 2 may set the count 20 of lost packets equal to zero. The threshold may be based, at least in part, upon a packet loss threshold and an average number of packets per frame.

The techniques described in this disclosure may be implemented within one or more of a general purpose microprocessor, digital signal processor (DSP), application specific integrated circuit (ASIC), field programmable gate array (FPGA), or other equivalent logic devices. Accordingly, the terms "processor" or "controller," as used herein, may refer to any one or more of the foregoing structures or any other structure suitable for implementation of the techniques described herein.

The various components illustrated herein may be realized by any suitable combination of hardware, software, firmware, or any combination thereof. In the figures, various 35 components are depicted as separate units or modules. However, all or several of the various components described with reference to these figures may be integrated into combined units or modules within common hardware and/or software. Accordingly, the representation of features as 40 components, units or modules is intended to highlight particular functional features for ease of illustration, and does not necessarily require realization of such features by separate hardware or software components. In some cases, various units may be implemented as programmable processes performed by one or more processors.

Any features described herein as modules, devices, or components may be implemented together in an integrated logic device or separately as discrete but interoperable logic devices. In various aspects, such components may be formed 50 at least in part as one or more integrated circuit devices, which may be referred to collectively as an integrated circuit device, such as an integrated circuit chip or chipset. Such circuitry may be provided in a single integrated circuit chip device or in multiple, interoperable integrated circuit chip 55 devices, and may be used in any of a variety of image, display, audio, or other multi-media applications and devices. In some aspects, for example, such components may form part of a mobile device, such as a wireless communication device handset.

If implemented in software, the techniques may be realized at least in part by a computer-readable data storage medium comprising code with instructions that, when executed by one or more processors, performs one or more of the methods described above. The computer-readable 65 medium may form part of a computer program product, which may include packaging materials. The computer-

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readable medium may comprise random access memory (RAM) such as synchronous dynamic random access memory (SDRAM), read-only memory (ROM), non-volatile random access memory (NVRAM), electrically erasable programmable read-only memory (EEPROM), embedded dynamic random access memory (eDRAM), static random access memory (SRAM), flash memory, magnetic or optical data storage media.

The techniques additionally, or alternatively, may be realized at least in part by a computer-readable communication medium that carries or communicates code in the form of instructions or data structures and that can be accessed, read, and/or executed by one or more processors. Any connection may be properly termed a computer-readable medium. For example, if the software is transmitted from a website, server, or other remote source using a coaxial cable, fiber optic cable, twisted pair, digital subscriber line (DSL), or wireless technologies such as infrared, radio, and microwave, then the coaxial cable, fiber optic cable, twisted pair, DSL, or wireless technologies such as infrared, radio, and microwave are included in the definition of medium. Combinations of the above should also be included within the scope of computer-readable media. Any software that is utilized may be executed by one or more processors, such as one or more DSP's, general purpose microprocessors, ASIC's, FPGA's, or other equivalent integrated or discrete logic circuitry.

Various aspects have been described in this disclosure. These and other aspects are within the scope of the following claims.

What is claimed is:

1. A method comprising:

detecting, by a video decoder, at least one error associated with coded video data sent from a video encoder;

determining a round-trip travel time for data communication between the video decoder and the video encoder:

determining a last-sent time at which the video decoder last sent the video encoder a previous feedback message indicating a previous error that was detected by the video decoder, wherein the previous error is a different error than the at least one error and is associated with different coded video data than the coded video data;

comparing a current time to a sum of the last-sent time and the round-trip travel time; and

based on the current time exceeding the sum of the last-sent time and the round trip travel time, sending the video encoder a feedback message indicating the at least one error detected by the video decoder, wherein the feedback message is different than the previous feedback message.

2. The method of claim 1, further comprising, upon sending the video encoder the feedback message,

setting the last-sent time equal to the current time.

- 3. The method of claim 1, wherein determining the round-trip travel time comprises measuring or estimating the round-trip travel time based upon at least one pair of messages that are exchanged between the video decoder and the video encoder.
  - **4**. The method of claim **1**, further comprising receiving the coded video data from the video encoder, and wherein detecting the at least one error comprises detecting at least one decoding error upon decoding of the coded video data by the video decoder.

- 5. The method of claim 1, wherein detecting the at least one error comprises detecting a packet loss, and wherein the method further comprises:
  - upon detecting the packet loss, incrementing a count of lost packets:
  - comparing the count of lost packets to a threshold;
  - based on the count of lost packets exceeding the threshold, sending the feedback message to the video encoder; and
  - based on the count of lost packets not exceeding the threshold, sending another type of feedback message indicating the packet loss to the video encoder.
  - 6. The method of claim 5, further comprising:
  - upon sending the feedback message to the video encoder, 15 setting the count of lost packets equal to zero.
- 7. The method of claim 5, wherein the threshold is based, at least in part, upon a packet loss threshold and an average number of packets per frame.
- **8**. The method of claim **4**, wherein detecting the at least 20 one error comprises detecting the at least one error in decoded video data from the video decoder.
  - 9. A device comprising:
  - a storage device storing coded video data sent from a video encoder; and
  - one or more processors coupled to the storage device and configured to:
    - detect at least one error associated with the coded video data sent from the video encoder;
    - determine a round-trip travel time for data communication between the device and the video encoder;
    - determine a last-sent time at which the device last sent the video encoder a previous feedback message indicating a previous error that was detected by the device, wherein the previous error is a different error than the at least one error and is associated with different coded video data than the coded video data:
    - compare a current time to a sum of the last-sent time and the round-trip travel time; and
    - based on the current time exceeding the sum of the last-sent time and the round trip travel time, send the video encoder a feedback message indicating the at least one error detected by the video decoder, wherein the feedback message is different than the 45 previous feedback message.
- 10. The device of claim 9, wherein the one or more processors are further configured to, upon sending the video encoder the feedback message, set the last-sent time equal to the current time.
- 11. The device of claim 9, wherein the one or more processors are configured to determine the round-trip travel time at least by measuring or estimating the round-trip travel time based upon at least one pair of messages that are exchanged between the device and the video encoder.
- 12. The device of claim 9, wherein the one or more processors are further configured to receive the coded video data from the video encoder, and wherein the one or more processors are configured to detect the at least one error by detecting at least one decoding error upon decoding of the coded video data by the device.
- 13. The device of claim 9, wherein the at least one error comprises a packet loss, and wherein the one or more processors are further configured to:
  - increment a count of lost packets upon detecting the packet loss;

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- compare the count of lost packets to a threshold;
- based on the count of lost packets exceeding the threshold, send the feedback message to the video encoder;
- based on the count of lost packets not exceeding the threshold, send another type of feedback message indicating the packet loss to the video encoder.
- 14. The device of claim 13, wherein the one or more processors are further configured to set the count of lost packets equal to zero upon sending the feedback message to the video encoder.
- 15. The device of claim 13, wherein the threshold is based, at least in part, upon a packet loss threshold and an average number of packets per frame.
- 16. The device of claim 9, wherein the device comprises a wireless communication device handset.
- 17. The device of claim 9, wherein the device comprises one or more integrated circuit devices.
- 18. The device of claim 12, wherein to detect the at least one error, the one or more processors are configured to detect the at least one error in decoded video data from the video decoder.
  - 19. A device comprising:
  - means for detecting at least one error associated with coded video data sent from a video encoder;
  - means for determining a round-trip travel time for data communication between the device and the video encoder;
  - means for determining a last-sent time at which the device last sent the video encoder a previous feedback message indicating a previous error that was detected by the device, wherein the previous error is a different error than the at least one error and is associated with different coded video data than the coded video data;
  - means for comparing a current time to a sum of the last-sent time and the round-trip travel time; and
  - means for, based on the current time exceeding the sum of the last-sent time and the round trip travel time, sending the video encoder a feedback message indicating the at least one error detected by the video decoder, wherein the feedback message is different than the previous feedback message.
  - 20. The device of claim 19, further comprising
  - means for, upon sending the video encoder the feedback message, setting the last-sent time equal to the current time.
- 21. The device of claim 19, wherein the means for determining the round-trip travel time comprises means for 50 measuring or estimating the round-trip travel time based upon at least one pair of messages that are exchanged between the video decoder and the video encoder.
  - 22. The device of claim 19, further comprising means for receiving the coded video data from the video encoder, and wherein the means for detecting the at least one error comprises means for detecting at least one decoding error upon decoding of the coded video data by the video decoder.
- 23. The device of claim 19, wherein the means for detecting the at least one error comprises means for detect-60 ing a packet loss, and wherein the device further comprises:
  - upon detecting the packet loss, means for incrementing a count of lost packets;
  - means for comparing the count of lost packets to a threshold:
  - means for, based on the count of lost packets exceeding the threshold, sending the feedback message to the video encoder; and

means for, based on the count of lost packets not exceeding the threshold, sending another type of feedback message indicating the packet loss to the video encoder.

24. The device of claim 23, further comprising:

upon sending the feedback message to the video encoder, 5 means for setting the count of lost packets equal to zero.

- 25. The device of claim 23, wherein the threshold is based, at least in part, upon a packet loss threshold and an average number of packets per frame.
- 26. The device of claim 22, wherein the means for detecting the at least one error comprises means for detecting the at least one error in decoded video data from the video decoder.
- 27. A non-transitory computer-readable storage medium 15 comprising instructions that, when executed, cause one or more processors to:

detect, by a video decoder, at least one error associated with coded video data sent from a video encoder;

determine a round-trip travel time for data communication 20 between the video decoder and the video encoder;

determine a last-sent time at which the video decoder last sent the video encoder a previous feedback message indicating a previous error that was detected by the video decoder, wherein the previous error is a different 25 error than the at least one error and is associated with different coded video data than the coded video data; compare a current time to a sum of the last-sent time and

compare a current time to a sum of the last-sent time and the round-trip travel time; and

based on the current time exceeding the sum of the 30 last-sent time and the round trip travel time, send the video encoder a feedback message indicating the at least one error detected by the video decoder, wherein the feedback message is different than the previous feedback message.

28. The non-transitory computer-readable storage medium of claim 27, wherein the instructions further cause the one or more processors to, upon sending the video encoder the feedback message

set the last-sent time equal to the current time.

29. The non-transitory computer-readable storage medium of claim 27, wherein instructions for causing the one or more processors to determine the round-trip travel

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time comprise instructions for causing the one or more processors to measure or estimate the round-trip travel time based upon at least one pair of messages that are exchanged between the video decoder and the video encoder.

- 30. The non-transitory computer-readable storage medium of claim 27, further comprising instructions for causing the one or more processors to receive the coded video data from the video encoder, and wherein instructions for causing the one or more processors to detect the at least one error comprise instructions for causing the one or more processors to detect at least one decoding error upon decoding of the coded video data by the video decoder.
- 31. The non-transitory computer-readable storage medium of claim 27, wherein instructions for causing the one or more processors to detect the at least one error comprise instructions for causing the one or more processors to detect a packet loss, and further comprising instructions for causing the one or more processors to:

upon detecting the packet loss, increment a count of lost packets;

compare the count of lost packets to a threshold;

based on the count of lost packets exceeding the threshold, send the feedback message to the video encoder; and

based on the count of lost packets not exceeding the threshold, send another type of feedback message indicating the packet loss to the video encoder.

- **32**. The non-transitory computer-readable storage medium of claim **31**, further comprising instructions for causing the one or more processors to, upon sending the feedback message to the video encoder, set the count of lost packets equal to zero.
- 33. The non-transitory computer-readable storage medium of claim 31, wherein the threshold is based, at least in part, upon a packet loss threshold and an average number of packets per frame.
  - **34.** The non-transitory computer-readable medium of claim **30**, wherein to detect the at least one error, the instructions cause the one or more processors to detect the at least one error in decoded video data from the video decoder.

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